

# Optimal Feedback in Contests

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with

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## Model (1/4): Players & Timing

- *Players*: A principal and  $n \geq 2$  agents
- At  $t = 0$ , the principal designs a contest comprising
  - i. a termination rule  $\tau$ ,
  - ii. a rule for allocating a \$1 prize, and
  - iii. a feedback policy.
- At every  $t > 0$ , each agent
  - receives a message per the feedback policy, and
  - chooses effort  $a_{i,t} \in [0, 1]$
- The contest ends at  $\tau$  and prize is awarded according to allocation rule

## Model (2/4): Agents' Output & "Who observes what"

- **Each agent's output** takes the form of a Poisson "breakthrough":
  - During  $(t, t + dt)$  agent  $i$  "succeeds" with probability  $a_{i,t}dt$
  - Each agent can succeed at most once
  - Denote  $x_{i,t} = 1$  if agent  $i$  has succeeded by  $t$ , and  $x_{i,t} = 0$  otherwise
- **Who observes what:**
  - Principal observes successes but not efforts
  - Each agent observes his effort but not successes
- Denote by  $p_{i,t}$  agent  $i$ 's belief at  $t$  that he has succeeded
  - *Note:* Effort is worthwhile for an agent only if he hasn't yet succeeded

## Model (3/4): Principal's Choice Variables

- A **termination rule**  $\tau$  is a stopping time w.r.t  $\mathbf{x}_t = \{x_{1,t}, \dots, x_{n,t}\}_{s \leq t}$ 
  - e.g., if  $\tau = \inf \{t : x_{i,t} = 1 \text{ some } i\}$ , contest ends upon first success
- A **prize allocation rule**  $\mathbf{q} \in [0, 1]^n$  specifies the probability each agent wins the prize as function of  $\mathbf{x}_\tau$ ; i.e., each agent's time of success
  - e.g., if  $q_i(\mathbf{x}_\tau) = \mathbb{I}_{\{x_{i,t} \geq x_{j,t} \forall j, t\}}$ , first agent to succeed wins prize w.p 1
- A **feedback policy**  $\mathcal{M}$  specifies the message sent to each agent at every  $t$  as a function of  $\mathbf{x}_t$  and past messages
  - e.g., if  $m_{i,t} = x_{i,t} \forall i, t$ , ea. agent is told whether he has succeeded
  - *Alternatives*: Random feedback, feedback about others' successes, feedback about feedback, etc

## Model (4/4): Payoffs

- Given a contest, each agent's expected utility is

$$u_{i,t} = \max_{a_i \in [0,1]} \mathbb{E} \left[ q_i(\mathbf{x}_\tau) - \int_t^\tau c a_{i,s} ds \right]$$

- First term*: Probability agent  $i$  wins the prize
  - Second term*: Cost of effort where  $c \in (1/n, 1)$
  - BNE*: Each agent chooses effort optimally anticipating rivals' efforts
- Principal chooses a contest  $\{\tau, \mathbf{q}, \mathcal{M}\}$  and effort recommendations to

$$\begin{aligned} \max_{\tau, \mathbf{q}, \mathcal{M}, \mathbf{a}} \mathbb{E} \left[ \sum_{i=1}^n \int_0^\tau a_{i,t} dt \right] \\ \text{s.t. } a_{i,t} \text{ is IC for all } i, t. \end{aligned}$$

## A Motivating Example

- Consider a manager who uses a promotion, acting as the prize, to motivate a group of employees
- Each agent must clear some “bar” to be eligible for promotion
  - This “bar” is represented by a success in the model (hence agents can succeed only once)
- Agents don’t definitively know whether they have cleared said “bar”, but the principal can disclose this (or other) information
- Manager cares about aggregate effort (not clearing the bar per-se)
- **Question:** How to design contest to get the most effort for \$1 prize?

# Remarks

## i. *No discounting.*

- Model is equivalent to one in which players discount time at some rate, and the value of the prize appreciates at the same rate

## ii. *Agents don't observe their own successes.*

- Goal is to give the principal full control of the agents' information
- In optimal contest, each agent is fully appraised of his own success; *i.e.*, main result would be unchanged if agents observed own successes

## iii. *Constant hazard rate of success.*

- Success during  $(t, t + dt)$  depends only on effort during this interval
- *Extension:* Arrival rate of success increases with past efforts

# Outline of Results

## **Proposition 1:** Optimal contest *without* feedback

- No messages permitted and contest ends at some deterministic  $T$
- *Egalitarian allocation rule* is optimal: Each agent who succeeds by  $T$  wins the prize with equal probability



## Outline of Results

### Proposition 2: Optimal contest *with* feedback

- *Cyclical structure:*
  - Initially, principal sets provisional deadline  $T^*$
  - If one or more agents succeed by  $T^*$ , contest ends at  $T^*$
  - Otherwise, the deadline is extended to  $t = 2T^*$
  - If no agent succeeds by  $2T^*$ , deadline extended until  $3T^*$ . And so on.
- When contest ends, prize is awarded according to egalitarian rule.
  - *i.e.*, every agent who succeeded wins prize with equal probability
- Agents are fully appraised of their own success. They are informed about their rivals' successes at  $T^*, 2T^*, \dots$ 
  - *i.e.*, if deadline is extended, then no one has succeeded yet
- This contest achieves the *first-best payoff* for the principal

# Outline of Results

**Proposition 3:** Optimal contest with increasing hazard rate

- Effort today makes success tomorrow more likely
- Similar structure, except that each provisional deadline has a stochastic duration

# No-feedback Contests

- First, we restrict attention to contests without feedback
  - No message transmission permitted (*i.e.*, no direct feedback)
  - Principal chooses a deterministic deadline  $T$  (*i.e.*, no indirect feedback)

- Fix a contest and for each agent, define **reward function**

$$R_{i,t} = \mathbb{E} [q_i(\mathbf{x}_T) \mid dx_{i,t} = 1]$$

*i.e.*, agent's expected reward conditional on succeeding at  $t$

- The agent's payoff can thus be expressed as

$$u_{i,t} = \max_{a_{i,s} \in [0,1]} \int_t^T (1 - p_{i,s}) a_{i,s} R_{i,s} - ca_{i,s} ds$$

- During  $(t, t + dt)$ , succeeds w.p  $(1 - p_{i,t})a_{i,t}dt$ , in which case earns  $R_{i,t}$ ,
- and he incurs cost  $ca_{i,t}dt$

# Agents' Problem

- Fix an arbitrary deadline  $T$  and reward function  $R_{i,t}$ . Agent solves

$$u_{i,0} = \max_{a_{i,t} \in [0,1]} \int_0^T [(1 - p_{i,t})R_{i,t} - c] a_{i,t} dt$$

s.t.  $\dot{p}_{i,t} = (1 - p_{i,t})a_{i,t}$  with  $p_{i,0} = 0$

- On the constraint:
  - Evolution equation for  $p_{i,t}$  follows from Bayes' rule
  - Captures fact that effort today lowers future probability of success
- Std. optimal control problem: Use Pontryagin's maximum principle

# Agents' Problem: Incentive Compatibility

- *Today's talk:* Restrict attention to contests with  $a_{i,t} = 1$  for all  $[0, T]$

## Lemma 1.

- Consider no-feedback contest w. deadline  $T$  and reward function  $R_{i,t}$
- Effort  $a_{i,t} = 1$  is incentive compatible for all  $t \in [0, T]$  if and only if

$$\underbrace{e^{-t} R_{i,t}}_{\text{MB at } t} \geq \underbrace{c}_{\text{direct MC}} + \underbrace{\int_t^T e^{-s} R_{i,s} ds}_{\text{strategic MC}} \text{ for all } t.$$

- 1<sup>st</sup> term: Success arrives at rate  $e^{-t}$ , and reward is  $R_{i,t}$
- 2<sup>nd</sup> term: (*Direct*) marginal cost of effort
- 3<sup>rd</sup> term: Success today eliminates possibility of success in the future

# No-feedback Contest: Principal's Problem

- Optimal no-feedback contest solves the following problem:

$$\max_{T, \mathbf{q}} n \int_0^T 1 dt$$

$$\text{s.t. } e^{-t} R_{i,t} \geq c + \int_t^T e^{-s} R_{i,s} ds \quad \forall i, t$$

$$T \geq 0, \mathbf{q} \text{ is a feasible prize allocation rule}$$

where  $R_{i,t} = \mathbb{E}[q_i(\mathbf{x}_\tau) \mid dx_{i,t} = 1]$ .

- The principal chooses
  - a terminal date  $T$ , and
  - a prize allocation rule  $\mathbf{q}$
 to maximize aggregate effort s.t IC constraint.
- Restriction to symmetric contests with max. effort shown to be wolog

# Optimal No-feedback Contest

- *Definition 1:* Egalitarian prize allocation rule (EGA)

$$\mathbf{q}_i^{ega}(\mathbf{x}_T) = \frac{x_{i,T}}{\sum_j x_{j,T}}$$

- *i.e.*, every agent who succeeds wins the prize with equal probability
- *Definition 2:*  $\hat{T}$  is the unique solution of  $1 - e^{-n\hat{T}} = nc(e^{\hat{T}} - 1)$ 
  - Given EGA & no feedback, this is longest max. effort is IC

## Proposition 1.

- The optimal no-feedback contest has deadline  $\hat{T}$  and egalitarian prize allocation rule  $\mathbf{q}^{ega}$ .
- In equilibrium, each agent exerts maximum effort for all  $t \in [0, \hat{T}]$ .

# Optimal No-feedback Contest: Heuristic Derivation (1/3)

- *Observation #1:*  $R_{i,t} = ce^T$  satisfies IC with equality for all  $t$ 
  - Time-invariant & symmetric  $R_{i,t}$  corresponds to EGA allocation rule
- *Observation #2:*
  - Recall  $R_{i,t}$  is prob. agent wins prize conditional on succeeding at  $t$
  - Given  $R_{i,t}$ , agent  $i$  wins the prize with probability  $\int_0^T e^{-t} R_{i,t} dt$ . So

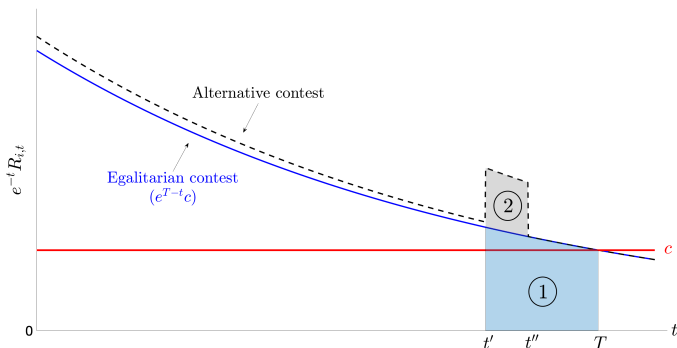
$$\underbrace{\sum_i \int_0^T e^{-t} R_{i,t} dt}_{\text{Pr}\{\text{prize is awarded}\}} \leq \underbrace{1 - e^{-nT}}_{\text{Pr}\{\text{at least one agent succeeds}\}}$$

- In other words, increasing  $e^{-t} R_{i,t}$  entails an opportunity cost, and so the principal wants to minimize  $e^{-t} R_{i,t}$  subject to satisfying IC.



# Optimal No-feedback Contest: Heuristic Derivation (2/3)

- Consider alternative contest with  $e^{-t}\tilde{R}_{i,t} > e^{-t}R_{i,t}$  on some interval



- Egalitarian contest*: IC at  $t'$  requires that  $e^{-t'}R_{i,t'} \geq c + \textcircled{1}$
- Alternative contest*: IC at  $t'$  requires that  $e^{-t'}\tilde{R}_{i,t'} \geq c + \textcircled{1} + \textcircled{2}$
- Thus  $e^{-t}\tilde{R}_{i,t} > e^{-t}R_{i,t}$  for all  $t < t'$ ; i.e.,  $\tilde{R}_{i,t}$  is more expensive

## Optimal No-feedback Contest: Heuristic Derivation (3/3)

- Thus, any non-egalitarian contest with deadline  $T$  can be replaced by EGA contest with same deadline that is *cheaper* for principal
  - *Cheaper*  $\Rightarrow$  Can extend deadline and still satisfy IC for all  $t$

- It remains to pin down the optimal deadline  $\hat{T}$ 
  - Fix a  $T$ . Given the egalitarian allocation rule,

$$\Pr \{ \text{agent } i \text{ wins prize} \} = \int_0^T e^{-t} R_{i,t}^{ega} dt = \frac{1 - e^{-nT}}{n}$$

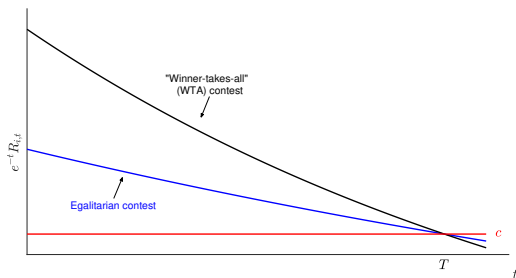
- Since  $R_{i,t}^{ega}$  is time-invariant, we have  $R_{i,t}^{ega} = [1 - e^{-nT}] / [n(1 - e^{-T})]$
- By def.  $\hat{T}$  is largest deadline for which  $R_{i,t}^{ega} \geq e^T c$ ; *i.e.*, max effort IC

## Optimality of Egalitarian Contest: Intuition (1/2)

- As an alternative, take “winner-takes-all” contest with deadline  $T$  *i.e.*, at  $T$ , the prize is awarded to the first agent who succeeded
- Assuming max. effort is IC on  $[0, T]$ , we have reward functions

$$R_{i,t}^{wta} = e^{-(n-1)t}$$

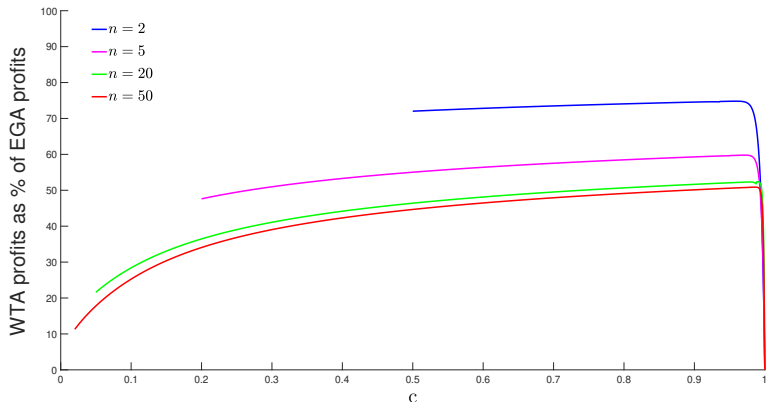
*i.e.*, if agent  $i$  succeeds at  $t$ , he is the first to do so w.p.  $e^{-(n-1)t}$



- Notice that  $e^{-t}R_{i,t}^{wta} > e^{-t}R_{i,t}^{ega}$ ; *i.e.*, WTA is more expensive than EGA

## Optimality of Egalitarian Contest: Intuition (2/2)

- The problem is that the WTA contest frontloads incentives too much
  - IC is slack for all  $t < T$ ; i.e., incentives excessively strong early on
- In contrast, EGA (maximally) backloads incentives s.t IC binds  $\forall t$



## Key Lemma: Sufficient Condition for Optimality

- Next, we consider contests with an arbitrary feedback policy

### Lemma 2:

A contest is guaranteed to be optimal if in equilibrium:

- The prize is awarded with probability 1; *i.e.*,  $\sum_i \mathbb{E}[q_i(\mathbf{x}_\tau)] = 1$
- Each agent earns zero rents; *i.e.*,  $u_{i,0} = 0$  for all  $i$

- The principal's object can be rewritten as

$$\mathbb{E} \left[ \sum_{i=1}^n \int_0^\tau a_{i,t} dt \right] = \frac{1}{c} \left[ \underbrace{\sum_i \mathbb{E}[q_i(\mathbf{x}_\tau)]}_{\text{Pr}\{\text{prize awarded}\} \leq 1} - \underbrace{\sum_i u_{i,0}}_{\text{rents} \geq 0} \right]$$

- If a contest attains those bounds, it must be optimal (and first-best)

## Step 1: Constructing a Zero-Rent Contest (1/2)

- We can write each agent's payoff as

$$u_{i,t} = \max_{a_{i,s} \in [0,1]} \int_t^\tau [(1 - p_{i,s})R_{i,s} - c] a_{i,s} ds$$

- For a contest to concede no rents to the agents,

$$(1 - p_{i,t})R_{i,t} = c \text{ for all } i, t$$

- Claim:* Whenever  $a_{i,t} > 0$ , such a contest must have  $p_{i,t} = 0$ 
  - Suppose there is an interval on which  $\dot{p}_{i,t} > 0$  and  $(1 - p_{i,t})R_{i,t} = c$
  - Agent can pause effort during first half of interval so  $p_{i,t}^{private} < p_{i,t}^{eqm}$
  - Then  $(1 - p_{i,t}^{private})R_{i,t} > c$ , so agent can earn rents during second half
- Thus feedback policy must keep agents appraised of own success
  - Define the feedback policy  $\mathcal{M}^{pronto} = \{m_{i,t} = x_{i,t} \text{ for all } i, t\}$

## Step 2: Constructing a Zero-Rent Contest (2/2)

- Since  $p_{i,t} = 0$  until each agent succeeds, contest must have

$$R_{i,t} = c \quad \text{for all } i, t$$

- For  $R_{i,t}$  to be time-invariant & symmetric, alloc. rule must be EGA
- Suppose prize is awarded according to EGA rule at some fixed  $T$
- My reward conditional on succeeding at  $t$ ,  $R_{i,t}$ , depends on how many rivals I expect to succeed by  $T$ 
  - This number  $N_T \sim \text{Binom}(n-1, 1 - e^{-T})$ , and

$$R_{i,t}^{ega} = \mathbb{E} \left[ \frac{1}{1 + N_T} \right]$$

- If  $T \simeq 0$ , no rivals will succeed a.s, so  $R_{i,t}^{ega} \simeq 1$
- As  $T \rightarrow \infty$ , all  $n-1$  of my rivals will succeed a.s, so  $R_{i,t}^{ega} \rightarrow 1/n$
- There is a unique  $T^*$  such that  $R_{i,t}^{ega} = c$

## Step 3: Towards an Optimal Contest

- Consider the contest with:
  - i. Deterministic deadline  $T^*$
  - ii. Egalitarian allocation rule
  - iii. Feedback policy  $\mathcal{M}^{pronto}$
- By construction,
  - $R_{i,t}^{ega} = c$  so ea agent exerts max. effort until he succeeds and  $u_{i,t} = 0$
  - But the prize is awarded with probability  $\sum_i \mathbb{E}[q_i(\mathbf{x}_\tau)] = 1 - e^{-nT^*} < 1$   
*i.e.*, this contest satisfies part (ii) of Lemma 2, but **not** part (i)
- Next, we amend this contest such that  $\sum_i \mathbb{E}[q_i(\mathbf{x}_\tau)] = 1$ 
  - By Lemma 2, such contest will be optimal.



## Step 3: Cyclical Structure

- Consider the (cyclical) termination rule:

$$\tau^* = \inf \left\{ t : t = kT^* , k \in \mathbb{N} \text{ and } \sum_i x_{i,t} \geq 1 \right\}$$

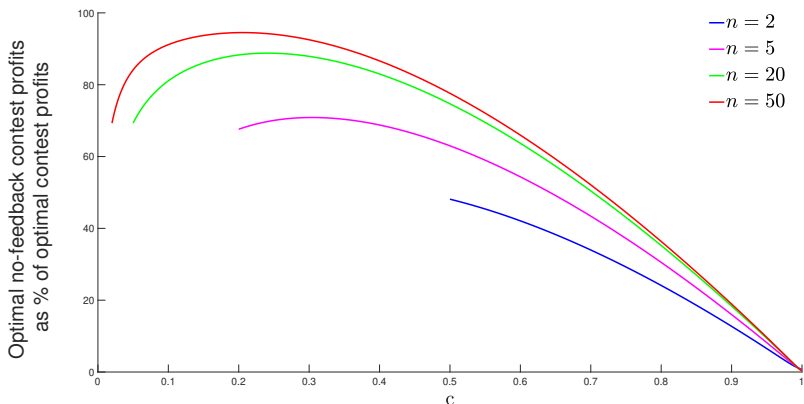
- This contest comprises “cycles” of length  $T^*$ , and is terminated at the end of the first cycle in which one or more agents have succeeded
- Within each cycle,  $R_{i,t}^{ega} = c$  by construction, so maximum effort is IC, and each agent’s instantaneous payoff is 0. Thus,  $u_{i,t} = 0$  for all  $t$ .
- Since the contest doesn’t end until at least one agent succeeds, the prize is awarded with probability 1.
  - i.e., the contest satisfies conditions of Lemma 2, and is hence optimal

# Optimal Contest (with Feedback)

## Proposition 2.

- The following contest is optimal:
  - i'. termination rule  $\tau^* = \inf \{t : t = kT^* , k \in \mathbb{N} \text{ and } \sum_i x_{i,t} \geq 1\}$ ,
  - ii. egalitarian prize allocation rule, and
  - iii. feedback policy  $\mathcal{M}^{pronto}$
- In equilibrium, each agent exerts max. effort until he succeeds
- Intuition for cyclical structure:
  - If rivals exert max. effort during  $[0, T]$ , my expected reward conditional on succeeding  $\downarrow T$  (because I will have to share prize with more rivals)
  - By construction,  $T^*$  is critical value such that  $R_{i,t} = c$
  - Cycles inform agents noone has succeeded, “resetting” incentives

# The Value of (optimal) Feedback



- Optimal feedback is most valuable when
  - Marginal costs  $c$  are small or large; *i.e.*, close to  $1/n$  or 1, or
  - Number of agents  $n$  is small

## Increasing Hazard Rate

- So far, we have assumed constant (unit) hazard rate of success
  - *i.e.*, agent succeeds during  $(t, t + dt)$  with probability  $a_{i,t}dt$
- Suppose instead that success arrives at rate  $\lambda_{i,t}a_{i,t}$ , and

$$\dot{\lambda}_{i,t} = f(\lambda_{i,t})a_{i,t}dt$$

for some function  $f(\cdot)$  and  $\lambda_{i,0} = \underline{\lambda}$ .

- I. Case  $f(\lambda) < 0$ : Effort today makes future success *less* likely
  - *e.g.*, Halac et al. (2017): “good news Poisson experimentation”
- II. Case  $f(\lambda) > 0$ : Effort today makes future success *more* likely
  - Optimal contest has similar features & properties as in base model: it awards the prize with probability 1 and extracts all rents

## Building Blocks

- Assume:  $f(\lambda) \geq 0$  and satisfies  $\lambda_{i,t} \in (c, nc)$ 
  - Suffices to assume  $\underline{\lambda} > c$  and  $f(\bar{\lambda}) = 0$  for some  $\bar{\lambda} \in (\underline{\lambda}, nc)$
- Let  $\lambda_t^*$  solve  $\dot{\lambda}_{i,t} = f(\lambda_{i,t})$  subject to  $\lambda_{i,t} = \underline{\lambda}$ 
  - This is the trajectory of  $\lambda_{i,t}$  if agent exerts max. effort
- By an earlier argument, feedback policy  $\mathcal{M}^{pronto}$  to extract all rents
- For max. effort to be IC and rents to be 0, we must have

$$\lambda_t^* R_{i,t} = c \quad \text{for all } i, t$$

- Because  $\lambda_t^*$  increases in  $t$ ,  $R_{i,t}$  must decrease in  $t$   
*i.e.*, incentives should be frontloaded since “earlier” success is “tougher”
- Suffices to find prize allocation and termination rules s.t  $R_{i,t} = c/\lambda_t^*$

# Optimal Contest

## Proposition 3.

- There exists an optimal contest from the following class:
  - ① *Cyclical stochastic structure*: Each cycle ends with rate  $\gamma(t, \lambda_t)$
  - ② At the end of each cycle, if a success has occurred, contest ends and prize is awarded according to EGA; otherwise, a new cycle starts
  - ③ Feedback policy  $\mathcal{M}^{pronto}$ ; *i.e.*, agents appraised of own success
- In equilibrium, each agent exerts max. effort until he succeeds
- *i.e.*, similar structure to before, except cycles have stochastic length
- If  $\gamma = \infty$ , contest is “winner-takes-all”
- If  $\gamma = 0$  for  $t < T$  and  $\gamma = \infty$  for  $t \geq T$ , the contest is egalitarian
- By choosing function  $\gamma(t, \lambda_t)$ , can fine-tune degree of frontloading

## Related Literature

- Static tournaments / contests:
  - Lazear & Rosen ('81), Green & Stokey ('83), Nalebuff & Stiglitz ('83)
  - *Optimal prize allocation*: Moldovanu & Sela ('01), Drugov & Ryvkin ('18, '19), Olszewski and Siegel ('20)
  - *"Turning down the heat"*: Fang et al. ('18) and Letina et al. ('20)
- Dynamic contests:
  - Taylor ('95), Benkert & Letina ('20)
  - *Tugs of war*: Moscarini & Smith ('11), Cao ('14)
- Feedback in contests:
  - *"Reveal intermediate progress?"*: Yildirim ('05), Lizzeri et al. ('05), Aoyagi ('10), Ederer ('10), Goltsman & Mukherjee ('19)
  - *Contests for experimentation*: Halac et al. ('17)

# Discussion

- Contest design with endogenous feedback
  - Cyclical structure
  - Egalitarian prize allocation rule (maximally backloads incentives)
  - Each agent is always appraised of own success, but is informed of rivals' successes periodically
- Future work
  - Continuous effort
  - Decreasing hazard rate
  - Continuous output / more general production functions
  - Asymmetric agents